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Topology-Free Querying of Protein Interaction Networks

Sharon Bruckner <u>Falk Hüffner</u> Richard M. Karp Ron Shamir Roded Sharan

Tel Aviv University

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Protein complexes

- A protein complex is a group of proteins which interact with each other to perform some task.
- Many protein complexes are known, in particular for model organisms like yeast.
- Problem: does a known protein complex also exists in the protein interaction network of another species?



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Complex query as Constrained Subgraph Isomorphism





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Complex query as CONSTRAINED SUBGRAPH ISOMORPHISM





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Complex query as Constrained Subgraph Isomorphism

Query



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Problems with Constrained Subgraph Isomorphism

• Not error tolerant



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- Not error tolerant
- Interactions between query proteins (*topology*) might not be available



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Problems with CONSTRAINED SUBGRAPH ISOMORPHISM

- Not error tolerant
- Interactions between query proteins (*topology*) might not be available
- Computationally very hard



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Complex Query as Colorful Connected Subgraph



COLORFUL CONNECTED SUBGRAPH

Input: An undirected, vertex colored graph *G*. **Output:** Find a connected subgraph of *G* whose vertices use each color exactly once (*colorful subgraph*).

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Dynamic Programming

Idea

Instead of looking at all $O(n^k)$ possible subgraphs, look only at $O(2^k)$ color sets.



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T[v, S] for $v \in V$ and S a set of colors: true if there is a connected subgraph of |S| vertices containing v with exactly the colors in S

$$T[v, S] = \bigvee_{\substack{u \in N(v) \\ S_1 \uplus S_2 = S}} T[v, S_1] + T[u, S_2] + w(u, v)$$



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Theorem

COLORFUL CONNECTED SUBGRAPH with k colors can be solved in $O(3^k|E|)$ time.

Fixed-parameter tractability

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Corollary

COLORFUL CONNECTED SUBGRAPH is fixed-parameter tractable with respect to k.

Integer Linear Programming

An Integer Linear Program (ILP) can maximize a linear function under linear constraints and integrality constraints.



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Binary variables

 $c_{v}, v \in V$: $v = 1 \iff v$ is part of the complex

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$$\gamma$$
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Central problem

Given a graph G = (V, E) and binary variables $c_v, v \in V$, find linear constraints such that $G[\{v \mid c_v = 1\}]$ is connected.

CONNECTED SUBGRAPH als ILP





CONNECTED SUBGRAPH als ILP





CONNECTED SUBGRAPH als ILP





CONNECTED SUBGRAPH als ILP



- An (arbitrary) selected vertex servers as sink
- Each other selected vertex is source of a flow of 1
- Only selected vertices take part in flow

Model extensions

• More than one color per vertex



Model extensions

- More than one color per vertex
- Insertions/Deletions



Model extensions

- More than one color per vertex
- Insertions/Deletions
- Maximize edge weight of complex



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Complete ILP

maximize
$$\sum_{(v,w)\in E} \omega_{vw} e_{vw}$$
 (1)

subject to

$$\sum_{v \in V} c_v = t \tag{2}$$

$$\sum_{v \in V} r_v = 1 \tag{3}$$

$$e_{vw} \le c_v \land e_{vw} \le c_w \qquad \forall (v,w) \in E \tag{4}$$

$$\begin{aligned} & \underset{w_{W}}{\overset{}} \geq 1/2c_{v} + 1/2c_{w} - 1/2 \quad \forall (v, w) \in E \\ & f_{ww} = -f_{wv} \qquad \forall (v, w) \in E \end{aligned} \tag{5}$$

$$\sum_{w \in N(v)} f_{vw} = c_v - tr_v \qquad \qquad \forall v \in V$$
(7)

$$f_{vw}, f_{wv} \leq (t-1)e_{vw} \qquad \forall (v,w) \in E$$
 (8)

$$\sum_{\nu \in V} g_{\nu\gamma} \le 1 \qquad \qquad \forall \gamma \in C \tag{10}$$

$$\sum_{v \in V} \sum_{\gamma \in \Gamma(v)} g_{v\gamma} = t - N_{\text{ins}}$$
(11

$$g_{v\gamma} \leq c_v \qquad \qquad \forall v \in V, \gamma \in \Gamma(v)$$

Bruckner et al. (TAU)

(12)



• First do data reduction



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 - If few colors, but large instance, use dynamic programming
 - Otherwise, use ILP





Protein-protein interaction networks:

- yeast (5 430 proteins, 39 936 interactions)
- fly (6650 proteins, 21275 interactions)
- human (7915 proteins, 28972 interactions)

Query several hundred complexes of size 4-25 from:

- yeast, fly, human (interaction information available)
- bovine, mouse, and rat (not enough interaction information available)



Algorithms

Experiments

Number of complexes found





Evaluation of results

- Functional coherence: Percentage of proposed complexes that are significantly enriched with "GO-Terms"
- Specificity: Percentage of proposed complexes that overlap significantly with known complexes



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Experiments

Quality of matches



TORQUE Web-Server

Input for query species

- Query complex
- (Enter a list of proteins or leave
- blank to use all FASTA file proteins)
- FASTA format sequences

Input for target species



O Upload my own target species data.

PPI network

FASTA format sequences

Set algorithm parameters

Interaction probability threshold [0.0-0.99] 0. BLAST threshold [1e-99..1e-3] 1E

http://www.cs.tau.ac.il/~bnet/torque.html

Bruckner et al. (TAU)



Saccharomyces cerevisiae 🗘

Browse...



TORQUE Web-Server



Blue: matched nodes in the target species. Within each node, top: target protein, bottom: the matching query protein. Grey: insertions of target proteins. The box lists the deleted query proteins, if any.

Best match for the DNA synthesome complex of the mouse in the network of yeast



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Topology-Free Querying of Protein Interaction Networks

Summary

- A topology-free querying model yields significant complex query results.
- With a combination of dynamic programming and ILP, even difficult instances can be solved optimally.

